

WHITTAKER FUNCTIONS AND DEMAZURE CHARACTERS

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With appendix by Dinakar Muthiah and Anna Puskás

Abstract In this paper, we consider how to express an Iwahori–Whittaker function through Demazure characters. Under some interesting combinatorial conditions, we obtain an explicit formula and thereby a generalization of the Casselman–Shalika formula. Under the same conditions, we compute the transition matrix between two natural bases for the space of Iwahori fixed vectors of an induced representation of a p -adic group; this corrects a result of Bump–Nakasuji.

Keywords: Iwahori–Whittaker functions; Casselman–Shalika formula; Demazure characters; Demazure–Lusztig operators; shellability

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1. Introduction

The *Casselman–Shalika formula* [6] describes a *spherical Whittaker function* using the root system and the character of an irreducible representation of the dual group. The formula not only plays a fundamental role in the theory of p -adic groups and automorphic forms, but also connects many different constructions in mathematics, such as Schubert varieties, crystal bases and Macdonald polynomials. For example, see [3].

In this paper, we study a generalization of the Casselman–Shalika formula to the case of *Iwahori–Whittaker functions* through *Demazure characters*. To be precise, let \mathfrak{g} be a finite-dimensional simple Lie algebra over \mathbb{C} , which should be considered as the Lie algebra of the dual group. Let P be the weight lattice of \mathfrak{g} , and $\mathbb{C}[P]$ the group algebra

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of P , with basis $e^\lambda, \lambda \in P$. The subset of dominant weights will be denoted by P_+ . We also denote by $\Phi \supset \Phi^+$ the set of roots and positive roots, by s_α the reflection corresponding to $\alpha \in \Phi$, by $\Pi = \{a_i\}_{i \in I}$ the set of simple roots, and by $S = \{\sigma_i\}_{i \in I}$ the set of simple reflections, which generates the Weyl group W . We denote reduced words by $s_1 \dots s_n$ with $s_i = s_{\alpha_i}$ for $\alpha_i \in \Pi$, which means that for each i we have $s_i = \sigma_j$ and $\alpha_i = a_j$ for some $j \in I$ (we use this notation in order to avoid double indexing in reduced words). Let v be an indeterminate, and set $\mathcal{O}_v = \mathbb{C}(v) \otimes \mathbb{C}[P]$.

Consider the Demazure character $\partial_{w,\lambda}$ for $w \in W$ and $\lambda \in P_+$, which is the formal character of the Demazure module associated with the weight $w\lambda$. When $w = w_\circ$, the longest element, the character $\partial_{w_\circ,\lambda}$ is nothing but the character of the irreducible representation of \mathfrak{g} with highest weight λ . Now the Casselman–Shalika formula [6] is given by

$$\tilde{W}_{w_\circ,\lambda} = \left(\prod_{\alpha \in \Phi^+} (1 - ve^{-\alpha}) \right) \partial_{w_\circ,\lambda}, \tag{1.1}$$

where $\tilde{W}_{w_\circ,\lambda}$ is the spherical Whittaker function.

As mentioned above, this paper is concerned with generalizing the formula (1.1) to the case involving the Iwahori–Whittaker functions $W_{w,\lambda}$ (to be defined in Section 2) and the Demazure characters $\partial_{x,\lambda}$, for $w, x \in W$. That is to say, we would like to compute the coefficients $C_{w,x} \in \mathcal{O}_v, x \leq w$, in the expansion

$$W_{w,\lambda} = \sum_{x \leq w} C_{w,x} \partial_{x,\lambda}.$$

To make the problem more tractable, we consider the *Demazure atoms* $D_{x,\lambda}$ (see §2), instead of working with the Demazure characters directly. We write

$$W_{w,\lambda} = \sum_{x \leq w} c_{w,x} D_{x,\lambda},$$

and study how to compute $c_{w,x} \in \mathcal{O}_v, x \leq w$. The coefficients $C_{w,x}$ and $c_{w,x}$ are related in a simple way (Corollary 2.2):

$$c_{w,x} = \sum_{x \leq y \leq w} C_{w,y} \quad \text{and} \quad C_{w,x} = \sum_{x \leq y \leq w} (-1)^{\ell(y) - \ell(x)} c_{w,y}.$$

Still, in general, it would be difficult to obtain a complete description of the coefficients $c_{w,x}$. However, the main result of this paper shows how to compute the coefficients $c_{w,x}$ under some interesting conditions involving *good words* and *shellability*. More precisely, under Condition (A) or (B) at the beginning of §5, we obtain the following theorem.

Theorem 1.1. *Let $w = s_1 \dots s_n$ be a reduced word with $s_i = s_{\alpha_i}$ for some $\alpha_i \in \Pi, i = 1, \dots, n$, and*

$$\beta_i = s_1 \dots \hat{s}_{i_1} \dots \hat{s}_{i_2} \dots \alpha_i, \quad i = 1, \dots, n,$$

where the indices $i_1 < \dots < i_d$ between 1 and n are determined by Condition (A) or (B). Then we have

$$c_{w,x} = (1 - ve^{-\beta_1}) \dots \mathcal{T}_{\beta_{i_1}} \left(\dots \mathcal{T}_{\beta_{i_d}} \left(\dots (1 - ve^{-\beta_n}) \dots \right) \right),$$

where $\mathcal{T}_\beta = (1 - ve^{-\beta})\partial_\beta - 1$ and ∂_β is the Demazure operator corresponding to the root β .

Conditions (A) and (B) are intriguing. In fact, based on thorough computer tests, in § 5.3 we conjecture that they are equivalent in a strong sense. Shortly after posting our paper, D. Muthiah and A. Puskás proved our conjecture; their proof is included as an Appendix. As discussed in § 4.1, Condition (A) is closely related to smoothness of Schubert varieties in flag varieties G/B . We also present some statistical information regarding the frequency with which these conditions are satisfied.

We establish an application of Conditions (A) and (B) to the problem of computing the transition matrix between two natural bases for the space of Iwahori fixed vectors of an induced representation of a p -adic group. The same problem was studied by Bump and Nakasuji [4]. Their result [4, Theorem 1.9] was that, when w admits a good word for x , the entry $m(x, w)$ of the transition matrix is given by

$$m(x, w) = \prod_{\alpha \in S(x, w)} \frac{1 - q^{-1} \mathbf{z}^\alpha}{1 - \mathbf{z}^\alpha}, \quad (1.2)$$

where $S(x, w)$ is the set of roots determined by the good word condition. However, the assumption of having a good word is not sufficient to deduce the factorization property. In fact, the statement and their proof of the above result are not valid, and there are counterexamples (see Remarks 5.1 and 6.3). More precisely, it is claimed below equation (5.3) in [4] that

$$s_1 \cdots \hat{s}_{j_1} \cdots \hat{s}_{j_k} \cdots s_{j-1} \hat{s}_j s_{j+1} \cdots \hat{s}_{i_m} \cdots s_n \leq s_1 \cdots \hat{s}_j \cdots s_n,$$

but this inequality does not follow from the good word condition. A similar inequality was claimed when Proposition 4.7 was applied in the same proof.

In § 6, we assume Condition (B) and prove the formula (1.2). The main idea of the proof is similar to that of [4]. In particular, the aforementioned inequalities are both guaranteed by Condition (B). Given Condition (A) or (B), which are actually equivalent (§ Appendix A), we ensure the factorization property and thus make a correction to the Bump–Nakasuji result in full root system generality. This provides another evidence that Conditions (A) and (B) are natural ones to be considered in representation theory.

Related to the above-mentioned coefficients $m(w, x)$, it is also worth noting the recent paper of Nakasuji and Naruse [10]. By using a change of basis in the Hecke algebra, they express all of these coefficients in a completely different way compared to (1.2), namely as sums over combinatorial sets. The mentioned change of basis in the Hecke algebra generalizes the theory of so-called *root polynomials*, which provides similar combinatorial formulas for localizations of Schubert classes in the equivariant cohomology and K -theory of flag varieties; see [9] and the references therein, as well as [10, Remark 1].

The fact that there are two types of formulas for the coefficients $m(w, x)$, namely the general formula in [10] and the simpler formula (1.2) if Condition (A) or (B) holds, is very similar to the existence of a general summation formula for Schubert classes (via root polynomials), versus a much simpler product formula in the smooth case; see [1, Chapter 7]. It turns out that the latter formula is hard to derive from the former, so completely separate proofs are needed. In this context, it is not surprising that Conditions (A) and (B) are related to smoothness of Schubert varieties, as noted above.

2. Description of the problem

In this section, we present the main question of this paper, introduced in Section 1, in more detail. We keep the notions fixed in Section 1.

Recall that the Hecke algebra \mathcal{H}_v is the algebra over $\mathbb{C}(v)$ defined by the generators T_i , $i \in I$, subject to the quadratic relations

$$T_i^2 = (v - 1)T_i + v, \quad i \in I,$$

and the braid relations corresponding to W . The algebra \mathcal{H}_v acts on \mathcal{O}_v by

$$T_i \mapsto \mathcal{T}_i := (1 - ve^{-a_i})\partial_i - 1, \quad i \in I,$$

where ∂_i , $i \in I$, are the Demazure operators defined by

$$\partial_i = \frac{1 - e^{-a_i}\sigma_i}{1 - e^{-a_i}}.$$

In particular, the operators \mathcal{T}_i , $i \in I$, which are known as the *Demazure–Lusztig operators*, satisfy the braid relations. Hence one may define

$$T_w \mapsto \mathcal{T}_w = \mathcal{T}_{i_1} \cdots \mathcal{T}_{i_l}$$

for an arbitrary choice of a reduced word $w = \sigma_{i_1} \cdots \sigma_{i_l}$. For a dominant weight $\lambda \in P_+$, define

$$W_{w,\lambda} = \mathcal{T}_w e^\lambda \quad \text{and} \quad \tilde{W}_{w,\lambda} = \sum_{x \leq w} W_{x,\lambda}, \quad w \in W.$$

As shown in [3], the expression $W_{w,\lambda}$ corresponds to the Iwahori–Whittaker function, and the sum $\tilde{W}_{w_\circ,\lambda}$ corresponds to the spherical Whittaker function where $w_\circ \in W$ is the longest element.

It is well known that the Demazure operators ∂_i , $i \in I$, satisfy the braid relations as well, so the operator ∂_w is well defined for $w \in W$ using any reduced word of w . Then the Demazure character is given by

$$\partial_{w,\lambda} = \partial_w e^\lambda, \quad \lambda \in P_+,$$

which is the formal character of the Demazure module associated with the weight $w\lambda$. Recall the Casselman–Shalika formula:

$$\tilde{W}_{w_\circ,\lambda} = \left(\prod_{\alpha \in \Phi_+} (1 - ve^{-\alpha}) \right) \partial_{w_\circ,\lambda}. \tag{2.1}$$

As mentioned above, we are interested in generalizing the formula (2.1) to the cases involving $\tilde{W}_{w,\lambda}$ (or $W_{w,\lambda}$) and $\partial_{x,\lambda}$ for $w, x \in W$. Precisely, we would like to compute the coefficients $\tilde{C}_{w,x} \in \mathcal{O}_v$, $x \leq w$, in the expansion

$$\tilde{W}_{w,\lambda} = \sum_{x \leq w} \tilde{C}_{w,x} \partial_{x,\lambda}.$$

Alternatively, if we write

$$W_{w,\lambda} = \sum_{x \leq w} C_{w,x} \partial_{x,\lambda},$$

we have

$$\tilde{C}_{w,x} = \sum_{x \leq y \leq w} C_{y,x}$$

and

$$C_{w,x} = \sum_{x \leq y \leq w} (-1)^{\ell(w)-\ell(y)} \tilde{C}_{y,x}$$

by the Möbius inversion [7, Theorem 1.2].

However, we found it more convenient to work with *Demazure atoms*. We define

$$D_i = \partial_i - 1 = e^{-a_i} \frac{1 - \sigma_i}{1 - e^{-a_i}}, \quad i \in I,$$

which is the specialization of \mathcal{T}_i at $v \rightarrow 0$. Then $D_i, i \in I$, satisfy the braid relations, and we define $D_w, w \in W$, in the obvious way. Now the Demazure atoms are defined to be

$$D_{w,\lambda} = D_w e^\lambda \quad \text{for } w \in W \text{ and } \lambda \in P_+.$$

Problem 1. Consider the transition between \mathcal{T}_w and D_w ,

$$\mathcal{T}_w = \sum_{x \leq w} c_{w,x} D_x,$$

and study how to compute $c_{w,x} \in \mathbf{Z}[v] \otimes \mathbf{Z}[P], x \leq w$.

The coefficients $C_{w,x}$ and $c_{w,x}$ can be related in a simple way, using the fact that the Demazure character is the sum of all the lower Demazure atoms. We give a proof of this fact below using a result in [3].

Lemma 2.1. $\partial_w = \sum_{x \leq w} D_x$ and $D_w = \sum_{x \leq w} (-1)^{\ell(w)-\ell(x)} \partial_x$.

Proof. $\partial_i, i \in I$, are the specialization of

$$\mathfrak{D}_i := \mathcal{T}_i + 1 = (1 - ve^{-a_i}) \partial_i$$

at $v \rightarrow 0$. Let \mathfrak{w} be a reduced word of w , and define $\mathfrak{D}_{\mathfrak{w}}$ in the obvious way. By [3, Theorem 6] which is proved using some results from [8], one has

$$\mathfrak{D}_{\mathfrak{w}} = \sum_{x \leq w} P_{\mathfrak{w},x}(v) \mathcal{T}_x,$$

where $P_{x,\mathfrak{w}}$ is the Poincaré polynomial of fiber of the Bott–Samelson resolution $Z_{\mathfrak{w}} \rightarrow X_w$ over the open cell $Y_x = Bx B/B$; the cohomology of this fiber was described combinatorially in [8]. Specializing $v \rightarrow 0$ gives that

$$\partial_w = \sum_{x \leq w} P_{x,\mathfrak{w}}(0) D_x = \sum_{x \leq w} D_x. \quad \square$$

Corollary 2.2. $c_{w,x} = \sum_{x \leq y \leq w} C_{w,y}$ and $C_{w,x} = \sum_{x \leq y \leq w} (-1)^{\ell(y)-\ell(x)} c_{w,y}$.

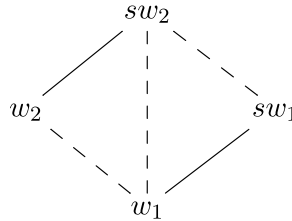


Figure 1. $Z(s, w_1, w_2)$ property.

By the reduction made above, the computation of the coefficients $\tilde{C}_{w,x}$ or $C_{w,x}$ is equivalent to the computation of the coefficients $c_{w,x}$, for $x \leq w$. Hence we will focus on Problem 1 from now on.

Note that the operators D_i are twisted derivations in the sense that

$$D_i(fg) = D_i(f) \cdot g + \sigma_i(f) \cdot D_i(g), \quad f, g \in \mathbf{Z}[P]. \tag{2.2}$$

In fact, the last equation is the specialization at $v \rightarrow 0$ of

$$\mathcal{T}_i(fg) = (1 - v)D_i(f) \cdot g + \sigma_i(f) \cdot \mathcal{T}_i(g), \quad f, g \in \mathbf{Z}[P]. \tag{2.3}$$

It is also known that $T_w, w \in W$, satisfy the relation

$$T_i \cdot T_w = \begin{cases} T_{\sigma_i w} & \text{if } \sigma_i w > w, \\ (v - 1)T_w + vT_{\sigma_i w} & \text{if } \sigma_i w < w. \end{cases} \tag{2.4}$$

For example, one has the quadratic relation $T_i^2 = (v - 1)T_i + v, i \in I$. Specializing (2.4) at $v \rightarrow 0$ gives

$$D_i \cdot D_w = \begin{cases} D_{\sigma_i w} & \text{if } \sigma_i w > w, \\ -D_w & \text{if } \sigma_i w < w. \end{cases} \tag{2.5}$$

3. Induction steps

In this section we give some general inductive steps for later use. We recall a well-known lemma from [7], which is called $Z(s, w_1, w_2)$ property of the Bruhat order, and it will be used frequently in this paper.

Lemma 3.1. *Let $s \in S$ be a simple reflection and $w_1, w_2 \in W$. Assume that $w_1 < sw_1, w_2 < sw_2$. Then*

$$w_1 \leq w_2 \iff w_1 \leq sw_2 \iff sw_1 \leq sw_2.$$

This lemma can be visualized using the diamond square in Figure 1, where the validities of the three dashed lines are all equivalent.

The following lemma can be easily verified by using (2.5).

Lemma 3.2. *Let $\alpha \in \Pi$ be a simple root and $s = s_\alpha$. Then*

$$\mathcal{T}_s \cdot D_w = \begin{cases} (1 - ve^{-\alpha})D_{sw} - ve^{-\alpha}D_w & \text{if } sw > w, \\ -D_w & \text{if } sw < w. \end{cases}$$

Lemma 3.3. *Assume that the simple reflection $s = s_\alpha$ is a left ascent of w , i.e., $sw > w$. Then*

$$\begin{aligned} \mathcal{T}_{sw} &= \sum_{x \leq w} \mathcal{T}_s(c_{w,x})D_x + \sum_{x \leq w, x < sx} (1 - ve^{-\alpha})s(c_{w,x})D_{sx} \\ &\quad - \sum_{x \leq w, x > sx} (1 - ve^{-\alpha})s(c_{w,x})D_x. \end{aligned}$$

Proof. Applying \mathcal{T}_s to the equation $\mathcal{T}_w = \sum_{x \leq w} c_{w,x}D_x$ and using (2.3) gives that

$$\mathcal{T}_{sw} = \sum_{x \leq w} s(c_{w,x})\mathcal{T}_s \cdot D_x + (1 - v)D_s(c_{w,x})D_x.$$

The lemma follows from inserting Lemma 3.2 into the last equation, and also from noting that

$$\begin{aligned} (1 - v)D_s - ve^{-\alpha}s &= (1 - ve^{-\alpha})\partial_s - 1 = \mathcal{T}_s, \\ (1 - v)D_s - s &= \mathcal{T}_s - (1 - ve^{-\alpha})s; \end{aligned}$$

here the second equation is an immediate consequence of the first. □

By comparing the coefficients in Lemma 3.3 with $\mathcal{T}_{sw} = \sum_{x \leq sw} c_{sw,x}D_x$, we obtain the following inductive algorithm.

Proposition 3.4. *Assume that $w < sw$, $s = s_\alpha \in S$, and that $x \leq sw$. Then*

(i) *if $x \leq w$, $x < sx$, then*

$$c_{sw,x} = \mathcal{T}_s(c_{w,x});$$

(ii) *if $x \leq w$, $x > sx$, then*

$$c_{sw,x} = (1 - ve^{-\alpha})s(c_{w,sx} - c_{w,x}) + \mathcal{T}_s(c_{w,x});$$

(iii) *if $x \not\leq w$, in which case $x > sx$, then*

$$c_{sw,x} = (1 - ve^{-\alpha})s(c_{w,sx}).$$

The three cases are illustrated in Figure 2. Note that in the last case we have either x and w incomparable, as depicted, or $x = sw > w = sx$.

The following corollary is immediate by applying Proposition 3.4(i) and (iii) recursively. Throughout, we let $\Phi_w := \Phi_+ \cap w\Phi_-$ be the inversion set of w^{-1} .

Corollary 3.5. *We have*

$$c_{w,e} = \mathcal{T}_w(1) \quad \text{and} \quad c_{w,w} = \prod_{\alpha \in \Phi_w} (1 - ve^{-\alpha}).$$

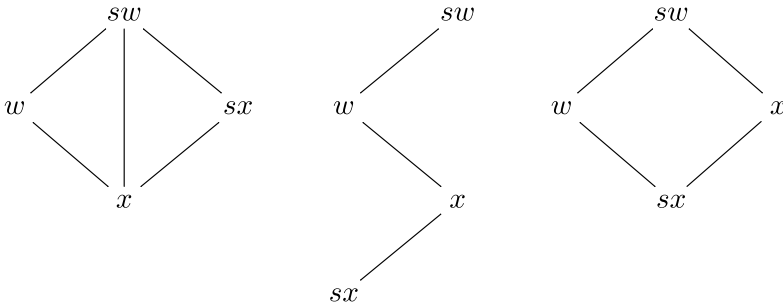


Figure 2. (i)–(iii) of Proposition 3.4.

4. Good words and shellability of Bruhat order

4.1. Good words

Following [4], we consider the notion of a good word. Assume that $x \leq w$, and introduce the sets

$$S(x, w) := \{\alpha \in \Phi_+ \mid x \leq ws_\alpha < w\}, \quad R(x, w) = \{s_\alpha \mid \alpha \in S(x, w)\}. \tag{4.1}$$

Deodhar’s inequality states that

$$\#S(x, w) = \#R(x, w) \geq \ell(w) - \ell(x), \tag{4.2}$$

with equality holding if the Kazhdan–Lusztig polynomial $P_{w \circ w, w \circ x} = 1$, or equivalently if the Schubert variety $X_{w \circ x}$ is rationally smooth at the T -fixed point $e_{w \circ w}$ (see [1]). We remark that $\#S(x, w)$ has the trivial upper bound $\ell(w)$ because of the inclusion $S(x, w) \subset \Phi_{w^{-1}} = \Phi_+ \cap w^{-1}\Phi_-$, where the last set is the inversion set of w , of cardinality $\ell(w)$; indeed, it is well known that $\alpha \in \Phi_+$ is an inversion of w , i.e., $w\alpha \in \Phi_-$, if and only if $ws_\alpha < w$.

For any reduced word $\mathbf{w} = s_1 \cdots s_n$ of w , let $\lambda_{x, \mathbf{w}}$ be the set of integers $i \in [1, n]$ such that $x \leq s_1 \cdots \hat{s}_i \cdots s_n$. Let $\alpha_i \in \Pi$ be such that $s_i = s_{\alpha_i}$, $i = 1, \dots, n$. Then there are bijections

$$\lambda_{x, \mathbf{w}} \rightarrow S(x, w) \rightarrow R(x, w), \quad i \mapsto \gamma_i := s_n \cdots s_{i+1}\alpha_i \mapsto s_{\gamma_i} = s_n \cdots s_{i+1}s_i s_{i+1} \cdots s_n.$$

Moreover it is clear that $ws_{\gamma_i} = s_1 \cdots \hat{s}_i \cdots s_n$. By abuse of notation, we also write

$$\lambda_{x, \mathbf{w}} = (i_1, \dots, i_d) \in \mathbf{N}^d \tag{4.3}$$

for the vector formed by elements of $\lambda_{x, \mathbf{w}}$ arranged in ascending order $i_1 < \cdots < i_d$. Then \mathbf{w} is called a *good word* for x if

$$x = s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_d} \cdots s_n. \tag{4.4}$$

Since $d = \#\lambda_{x, \mathbf{w}} \geq \ell(w) - \ell(x)$, a good word exists only if (4.4) is a reduced word; hence $d = \ell(w) - \ell(x)$. Conversely, it is conjectured in [4] that if W is simply-laced and $d = \ell(w) - \ell(x)$, then w has a good word for x . This conjecture is proved in [4] for $W = A_4$ or D_4 using SAGE, and it is shown to be false in the non-simply-laced case, e.g., for $W = B_2$.

4.2. Shellability

We recall the lexicographic shellability of Bruhat order, following [2]. For $x, y \in W$, we say that y covers x , denoted by $y \rightarrow x$, if $y > x$ and there is no $z \in W$ such that $y > z > x$. In this case $\ell(y) = \ell(x) + 1$ and there is a unique $\alpha \in \Phi_+$ such that $s_\alpha y = x$. Moreover for any reduced word $y = s_1 \dots s_l$, there is a unique $1 \leq i \leq l$ such that $x = s_1 \dots \hat{s}_i \dots s_l$, and one has $\alpha = s_1 \dots s_{i-1} \alpha_i$. We may also write $y \xrightarrow{\alpha} x$ to specify the reflection s_α that takes y to x .

Consider $x \leq w$ and the Bruhat interval $[x, w] := \{y \in W | x \leq y \leq w\}$. Then all maximal chains $\mathcal{C} : w = w_0 \rightarrow w_1 \rightarrow \dots \rightarrow w_d = x$ of $[x, w]$ have the same length $d = \ell(w) - \ell(x)$. Let us describe a labeling of the maximal chains of $[x, w]$. Fix once for all a reduced word $\mathbf{w} = s_1 \dots s_n$ of w . For a maximal chain \mathcal{C} of $[x, w]$ as above, there is a unique sequence i_1, \dots, i_d of distinct integers in $[1, n]$ such that w_k is obtained by removing s_{i_1}, \dots, s_{i_k} from \mathbf{w} , $k = 1, \dots, d$. In particular, this implies that the resulting subwords representing w_k 's are all reduced. Then we assign the label

$$\lambda(\mathcal{C}) = (\lambda_1(\mathcal{C}), \dots, \lambda_d(\mathcal{C})) := (i_1, \dots, i_d) \in \mathbf{N}^d. \tag{4.5}$$

Recall that the lexicographic order of \mathbf{N}^d is the linear ordering $<_L$ such that $\mathbf{a} = (a_1, \dots, a_d) <_L \mathbf{b} = (b_1, \dots, b_d)$ if $a_i < b_i$ in the first coordinate where they differ. The main result of [2] states that $[x, w]$ is lexicographically shellable. In particular, this implies that:

- (i) there is a unique maximal chain $\mathcal{C}_{x, \mathbf{w}}^+$ in $[x, w]$ whose labeling $\lambda(\mathcal{C}_{x, \mathbf{w}}^+)$ is increasing, i.e., $\lambda_1(\mathcal{C}_{x, \mathbf{w}}^+) < \dots < \lambda_d(\mathcal{C}_{x, \mathbf{w}}^+)$;
- (ii) $\lambda(\mathcal{C}_{x, \mathbf{w}}^+) <_L \lambda(\mathcal{C})$ for any other maximal chain \mathcal{C} of $[x, w]$.

Note that the maximal chain $\mathcal{C}_{x, \mathbf{w}}^+$ depends on the choice of the reduced word \mathbf{w} , which we fix from the beginning.

Similarly, consider the reduced word $s_n \dots s_1$ of w^{-1} . By applying shellability to w^{-1} with this reduced word and reverting to w , we see that:

- (i') there is a unique maximal chain $\mathcal{C}_{x, \mathbf{w}}^-$ in $[x, w]$ whose label $\lambda(\mathcal{C}_{x, \mathbf{w}}^-)$ is decreasing, i.e., $\lambda_1(\mathcal{C}_{x, \mathbf{w}}^-) > \dots > \lambda_d(\mathcal{C}_{x, \mathbf{w}}^-)$;
- (ii') $\lambda(\mathcal{C}_{x, \mathbf{w}}^-) >_L \lambda(\mathcal{C})$ for any other maximal chain \mathcal{C} in $[x, w]$.

5. Main result

In this section we compute the coefficient $c_{w,x}$, for $x \leq w$, under either of the following two conditions for the pair (w, x) :

- (A) w admits a reduced word \mathbf{w} such that $\lambda_{x, \mathbf{w}} = \lambda(\mathcal{C}_{x, \mathbf{w}}^-)^* = (i_1, \dots, i_d)$;
- (B) w admits a reduced word \mathbf{w} such that $\lambda(\mathcal{C}_{x, \mathbf{w}}^+) = \lambda(\mathcal{C}_{x, \mathbf{w}}^-)^* = (i_1, \dots, i_d)$.

Here we write $\lambda^* = (i_d, \dots, i_1) \in \mathbf{N}^d$ for a vector $\lambda = (i_1, \dots, i_d) \in \mathbf{N}^d$. Note that the reduced word \mathbf{w} satisfying Condition (A) is a good word for x .

Remark 5.1. Conditions (A) and (B) are stronger than the existence of a good word for x . Indeed, in type G_2 consider w with the (unique) reduced word $\mathbf{w} = \sigma_1 \sigma_2 \sigma_1 \sigma_2 \sigma_1$

and $x = \sigma_1$. We have

$$\lambda_{x, \mathfrak{w}} = (1, 2, 4, 5), \quad \lambda(\mathcal{C}_{x, \mathfrak{w}}^-)^* = (2, 3, 4, 5), \quad \lambda(\mathcal{C}_{x, \mathfrak{w}}^+) = (1, 2, 3, 4);$$

so \mathfrak{w} is a good word for x , but Conditions (A) and (B) fail. In fact, such examples can be found in simply-laced types too, for instance in type D_4 consider the reduced word $\mathfrak{w} = \sigma_2\sigma_1\sigma_3\sigma_4\sigma_2\sigma_4\sigma_3\sigma_1\sigma_2$ and $x = \sigma_2$ (where 2 is the central node of the Dynkin diagram). We have

$$\lambda_{x, \mathfrak{w}} = (1, 2, 3, 4, 6, 7, 8, 9), \quad \lambda(\mathcal{C}_{x, \mathfrak{w}}^-)^* = (2, 3, 4, 5, 6, 7, 8, 9),$$

$$\lambda(\mathcal{C}_{x, \mathfrak{w}}^+) = (1, 2, 3, 4, 5, 6, 7, 8).$$

It is easy to see that the same holds for any reduced word for the given Weyl group element. So again the above \mathfrak{w} is a good word for x , but Conditions (A) and (B) fail.

As we will prove, both Conditions (A) and (B) guarantee that only the relations in Proposition 3.4(i) and (iii) are used in the recursive computation of $c_{w,x}$; these relations have the advantage of being simple, compared to the relation in part (ii).

5.1. Lemmas on good words and shellability

We first prove a few more facts regarding combinatorial properties of a reduced word.

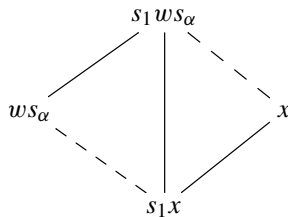
Lemma 5.2. *Assume that $\mathfrak{w} = s_1 \cdots s_n$ is a good word of w for x such that $\lambda_{x, \mathfrak{w}} = (i_1, \dots, i_d)$ with $i_1 > 1$. Then*

- (i) $x \not\leq s_1 w$;
- (ii) $S(s_1 x, s_1 w) = S(x, w)$;
- (iii) $s_1 \mathfrak{w} := s_2 \cdots s_n$ is a good word of $s_1 w$ for $s_1 x$ and $\lambda_{s_1 x, s_1 \mathfrak{w}} = (i_1 - 1, \dots, i_d - 1)$.

Proof. Part (i) is obvious from the definition of good word. Part (iii) follows from (ii). To prove (ii), it suffices to show that $S(s_1 x, s_1 w)$ is contained in $S(x, w)$, which implies that $S(s_1 x, s_1 w) = S(x, w)$ because of Deodhar’s inequality

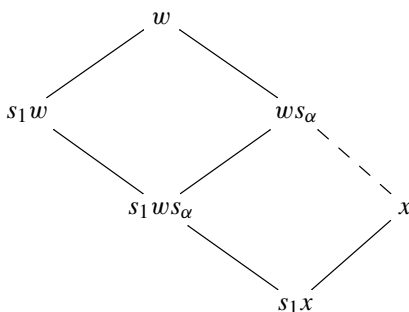
$$\#S(s_1 x, s_1 w) \geq \ell(s_1 w) - \ell(s_1 x) = \ell(w) - \ell(x) = \#S(x, w).$$

Take $\alpha \in S(s_1 x, s_1 w)$, i.e., $s_1 x \leq s_1 w s_\alpha < s_1 w$. We claim that $s_1 w s_\alpha < w s_\alpha$. To the contrary, assume that $s_1 w s_\alpha > w s_\alpha$. Then by Lemma 3.1 we have the diamond square



where the two dashed lines follow from the middle vertical line. This implies that $x \leq s_1 w s_\alpha < s_1 w$, a contradiction to part (i). Hence $s_1 w s_\alpha < w s_\alpha$, and using Lemma 3.1

again we obtain the diagram



which implies that $\alpha \in S(x, w)$. □

Lemma 5.3. *Let $w = s_1 \cdots s_n$ be a fixed reduced word of w , $\lambda(\mathcal{C}_{x,w}^+) = (i_1, \dots, i_d)$, $\lambda(\mathcal{C}_{x,w}^-) = (j_d, \dots, j_1)$, where $i_1 < \dots < i_d$ and $j_1 < \dots < j_d$. Consider the reduced word $s_1 w = s_2 \cdots s_n$ of $s_1 w$. Then*

(i) *if $i_1 > 1$, then $x \not\leq s_1 w$ and*

$$\lambda(\mathcal{C}_{s_1 x, s_1 w}^+) = (i_1 - 1, \dots, i_d - 1);$$

(ii) *if $j_1 > 1$, then*

$$\lambda(\mathcal{C}_{s_1 x, s_1 w}^-) = (j_d - 1, \dots, j_1 - 1);$$

(iii) *if $i_1 = 1$, then*

$$\lambda(\mathcal{C}_{x, s_1 w}^+) = (i_2 - 1, \dots, i_d - 1);$$

(iv) *if $j_1 = 1$, then $x < s_1 x$ and*

$$\lambda(\mathcal{C}_{x, s_1 w}^-) = (j_d - 1, \dots, j_2 - 1).$$

Proof. Write $\mathcal{C}_{x,w}^\pm : w = w_0^\pm \rightarrow w_1^\pm \rightarrow \dots \rightarrow w_d^\pm = x$.

(i) The last claim is clear since we have obviously a maximal chain

$$\mathcal{C} : s_1 w = s_1 w_0^+ \rightarrow s_1 w_1^+ \rightarrow \dots \rightarrow s_1 w_d^+ = s_1 x$$

of $[s_1 x, s_1 w]$ with increasing label $\lambda(\mathcal{C}) = (i_1 - 1, \dots, i_d - 1)$. We must have $\mathcal{C} = \mathcal{C}_{s_1 x, s_1 w}^+$ because of the uniqueness of the increasing label. It remains to prove that $x \not\leq s_1 w$. To the contrary, assume that $x \leq s_1 w = s_2 \cdots s_n$. Then concatenation of $w \rightarrow s_1 w$ with any maximal chain in $[x, s_1 w]$ will give a maximal chain \mathcal{C} in $[x, w]$ such that $\mathcal{C} <_L \mathcal{C}_{x,w}^+$, since $\lambda_1(\mathcal{C}) = 1 < \lambda_1(\mathcal{C}_{x,w}^+) = i_1$. This is a contradiction.

(ii) The proof is similar.

(iii) $\mathcal{C}_{x, s_1 w}^+$ equals the following subchain of $\mathcal{C}_{x,w}^+$

$$s_1 w = w_1^+ \rightarrow w_2^+ \rightarrow \dots \rightarrow w_d^+ = x.$$

(iv) $s_1x \rightarrow x$ is the last arrow in the chain $\mathcal{C}_{x,w}^-$, and hence $x < s_1x$. The following subchain of $\mathcal{C}_{x,w}^-$

$$w = w_0^- \rightarrow w_1^- \rightarrow \dots \rightarrow w_{d-1}^- = s_1x$$

gives rise to the maximal chain of $[x, s_1w]$

$$\mathcal{C} : s_1w = s_1w_0^- \rightarrow s_1w_1^- \rightarrow \dots \rightarrow s_1w_{d-1}^- = x$$

with decreasing label $\lambda(\mathcal{C}) = (j_d - 1, \dots, j_2 - 1)$, which implies that $\mathcal{C} = \mathcal{C}_{x,s_1w}^-$. □

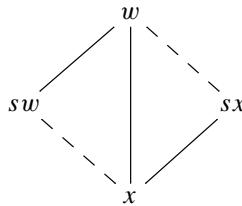
Lemma 5.4. *Assume that $w = s_1 \dots s_n$ is a good word of w for x such that $\lambda_{x,w} = \lambda(\mathcal{C}_{x,w}^-)^* = (i_1, \dots, i_d)$ with $i_1 = 1$. Then*

- (i) $x < s_1x$;
- (ii) $S(x, s_1w) = S(x, w) \setminus \{\gamma_1\}$, where $\gamma_1 = s_n \dots s_2\alpha_1$;
- (iii) $s_1w = s_2 \dots s_n$ is a good word of s_1w for x ;
- (iv) $\lambda_{x,s_1w} = (i_2 - 1, \dots, i_d - 1) = \lambda(\mathcal{C}_{x,s_1w}^-)^*$.

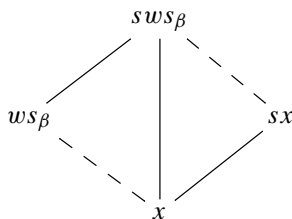
Proof. Part (i) and the last equality in (iv) follow from Lemma 5.3(iv). Part (iii) and the first equality in (iv) are direct consequences of (ii). Finally (ii) follows from Lemma 5.5 below, which is of independent interest. □

Lemma 5.5. *Assume that $x < w$, $sw < w$ and $x < sx$, where $s = s_\alpha \in S$. If $\#S(x, w) = \ell(w) - \ell(x)$, then $S(x, sw) = S(x, w) \setminus \{-w^{-1}\alpha\}$.*

Proof. Consider the following diamond given by Lemma 3.1.



Take $\beta \in S(x, sw)$, i.e., $sw > sws_\beta \geq x$. We claim that $\beta \in S(x, w)$, i.e., $w > ws_\beta \geq x$. If $ws_\beta > sws_\beta$, then the claim is obvious, again by Lemma 3.1. If $ws_\beta < sws_\beta$, then Lemma 3.1 gives the following diamond



Hence the claim follows. Obviously $\beta \neq -w^{-1}\alpha \in S(x, w)$, because $sws_{w^{-1}\alpha} = w > sw$. Therefore we get an inclusion $S(x, sw) \subset S(x, w) \setminus \{-w^{-1}\alpha\}$. This inclusion is an equality

because of Deodhar’s inequality

$$\#S(x, sw) \geq \ell(sw) - \ell(x) = \ell(w) - \ell(x) - 1 = \#S(x, w) - 1,$$

where the last equality follows from the assumption $\#S(x, w) = \ell(w) - \ell(x)$. □

5.2. Main theorem

We can now apply previous lemmas together with Proposition 3.4 recursively to compute $c_{w,x}$, assuming Condition (A) or (B). As mentioned above, only cases (i) and (iii) of Proposition 3.4 show up in the computation. In order to formulate our main result, we introduce an additional notation.

For any $\alpha \in \Phi$, let

$$\partial_\alpha = \frac{1 - e^{-\alpha} s_\alpha}{1 - e^{-\alpha}}, \quad \mathcal{T}_\alpha = (1 - ve^{-\alpha})\partial_\alpha - 1. \tag{5.1}$$

Using this notation, it is easy to see that we have

$$w \cdot \partial_\alpha = \partial_{w\alpha} \cdot w, \quad w \cdot \mathcal{T}_\alpha = \mathcal{T}_{w\alpha} \cdot w. \tag{5.2}$$

Theorem 5.6. *Assume that either condition (A) or (B) holds. In either case, let*

$$\beta_i = s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_2} \cdots \alpha_i, \quad i = 1, \dots, n.$$

Then we have

$$c_{w,x} = (1 - ve^{-\beta_1}) \cdots \mathcal{T}_{\beta_{i_1}} \left(\cdots \mathcal{T}_{\beta_{i_d}} \left(\cdots (1 - ve^{-\beta_n}) \right) \cdots \right).$$

Proof. In either case we use recursion. First assume Condition (A). If $i_1 > 1$, then $(s_1 w, s_1 x)$ satisfies Condition (A) as well, due to Lemma 5.2(iii) and Lemma 5.3(ii). Moreover, we may apply Proposition 3.4(iii) because of Lemma 5.2(i), which gives that

$$c_{w,x} = (1 - ve^{-\alpha_1}) s_1(c_{s_1 w, s_1 x}).$$

If $i_1 = 1$, then $(s_1 w, x)$ also satisfies (A) and we may apply Proposition 3.4(i), due to Lemma 5.4, which gives that

$$c_{w,x} = \mathcal{T}_{\alpha_1}(c_{s_1 w, x}).$$

Iterating this process gives us

$$c_{w,x} = (1 - ve^{-\alpha_1}) s_1 \cdots \mathcal{T}_{\alpha_{i_1}} \left(\cdots \mathcal{T}_{\alpha_{i_d}} \left(\cdots (1 - ve^{-\alpha_n}) s_n(1) \right) \cdots \right).$$

One may use (5.2) to push the reflections s_i , $1 \leq i \leq n$, $i \neq i_1, \dots, i_d$, across the operators $\mathcal{T}_{\alpha_{i_1}}, \dots, \mathcal{T}_{\alpha_{i_d}}$, noting that $x(1) = 1$.

The proof assuming Condition (B) is similar. If $i_1 > 1$, then by Lemma 5.3(i)–(ii), $(s_1 w, s_1 x)$ also satisfies (B) and Proposition 3.4(iii) applies. If $i_1 = 1$, then by Lemma 5.3(iii)–(iv), $(s_1 w, x)$ satisfies (B) and Proposition 3.4(i) applies. □

Remarks 5.7. (i) Note that in the special cases $d = n$ and $d = 0$, we recover $c_{w,e}$ and $c_{w,w}$ respectively, as given by Corollary 3.5.

(ii) The roots β_i can be interpreted as follows. We have

$$\{\beta_i : 1 \leq i \leq n, i \neq i_1, \dots, i_d\} = \Phi_x = \Phi_+ \cap x\Phi_-.$$

Moreover, under Condition (B), the roots $\beta_{i_1}, \dots, \beta_{i_d}$ give the sequence of reflections along the maximal chain $\mathcal{C}_{x,w}^+$ of $[x, w]$, i.e., we have

$$\mathcal{C}_{x,w}^+ : w = w_0 \xrightarrow{\beta_{i_1}} w_1 \rightarrow \dots \xrightarrow{\beta_{i_d}} w_d = x.$$

Hence the calculation of $c_{w,x}$ amounts to inserting the operators

$$\mathcal{T}_\beta = (1 - ve^{-\beta})\partial_\beta - 1, \quad \beta \in \{\beta_{i_1}, \dots, \beta_{i_d}\}$$

into the product $\prod_{\alpha \in \Phi_x} (1 - ve^{-\alpha})$ in a natural, combinatorial way.

5.3. Conditions (A) and (B)

Since these conditions are essential for our main result, we now discuss them in more detail. We start with an example.

Example 5.8. Consider $w = s_1s_2s_1s_3s_2s_1$ and $x = s_2s_3$ in A_3 . It is easy to see that both Conditions (A) and (B) hold in this example.

Based on thorough computer tests, we now formulate a conjecture about the equivalence of Conditions (A) and (B) in a strong sense.

Conjecture 5.9. *Let w be a reduced word for w and $x \leq w$. The following are equivalent:*

- (i) $\lambda_{x,w} = \lambda(\mathcal{C}_{x,w}^-)^*$;
- (ii) $\lambda(\mathcal{C}_{x,w}^+) = \lambda(\mathcal{C}_{x,w}^-)^*$;
- (iii) $\lambda_{x,w} = \lambda(\mathcal{C}_{x,w}^+)$.

The proof of the conjecture, due to D. Muthiah and A. Puskás, is included as an Appendix. Thus, we are able use the more symmetric condition

$$\lambda_{x,w} = \lambda(\mathcal{C}_{x,w}^+) = \lambda(\mathcal{C}_{x,w}^-)^*.$$

Note that it is enough to prove (i) \Leftrightarrow (ii), as (i) \Leftrightarrow (iii) would easily follow; indeed, just reverse the reduced word and use the fact that inversion is an automorphism of the Bruhat order.

We now discuss some statistics related to the frequency with which Conditions (A) and (B) are satisfied. We looked at the symmetric groups S_4, S_5 , and S_6 , as well as at the hyperoctahedral groups B_4 and B_5 . For each (signed) permutation w , we calculated the percentage of $x \leq w$, which satisfy Conditions (A) and (B). (The computer calculation is based on the following idea: for each w and reduced word for it, we generated via depth-first search all saturated decreasing chains in Bruhat order with increasing/decreasing labelings, which produces the corresponding x ; then, given this data for each such x , Conditions (A) and (B) are easily tested.) The distribution of

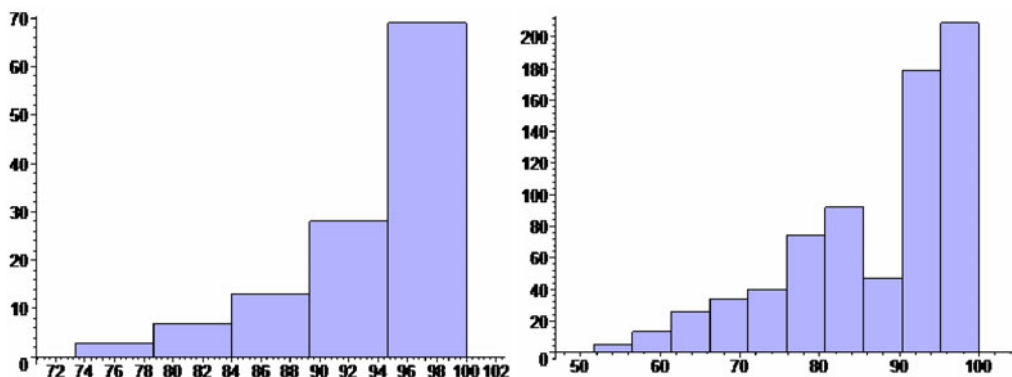


Figure 3. Histograms for S_5 and S_6 .

the mentioned percentages in S_5 and S_6 is shown in Figure 3, where the vertical axis represents the number of permutations. It is interesting to note that this distribution is skewed right, with the mode at the right tail, while the interquartile range reaches 100% in both cases. By contrast, in type B , the distribution looks closer to a uniform one.

Experiments with the same Weyl groups mentioned above also showed that the formula in Theorem 5.6 fails if Conditions (A) and (B) are not satisfied.

6. Casselman’s basis of Iwahori vectors

In this section, under the shellability Condition (B) in Section 6, we compute the transition matrix between two natural bases of the Iwahori fixed vectors in a spherical representation of a semisimple p -adic group, considered by Casselman in [5]. A conjectural formula is given in [4], which is proved under the assumption that a good word exists; however, it seems that there is a gap in this proof, which we do not know how to fix at present. We follow the strategy of computations in [4], although we consider reduced words from a very different point of view. Let us first recall the basic formulations and collect a few results we need from [4].

Let $\chi = \chi_{\mathbf{z}}$ be an unramified character of $T(F)$, which is parametrized by an element \mathbf{z} in the complex torus \hat{T} of the L-group ${}^L G$. Let $V(\chi) = \text{Ind}_B^G \chi$ be the induced representation, which consists of locally constant functions $f : G \rightarrow \mathbf{C}$ such that $f(bg) = (\delta^{1/2} \chi)(b)f(g)$, where $\delta = \det(\text{Ad}|_{\mathfrak{n}})$ is the modular character. Let J be the Iwahori subgroup, which is the preimage of $B(\mathbf{F}_q)$ under the reduction $K = G(\mathcal{O}_F) \rightarrow G(\mathbf{F}_q)$. Then the space of J -fixed vectors $V(\chi)^J$ has dimension $|W|$, and there are two bases $\{\phi_{w,\chi}\}$ and $\{f_w\}$ of $V(\chi)^J$ parametrized by W .

The first natural basis $\{\phi_{w,\chi}\}$ is defined using the disjoint decomposition $G = \bigsqcup_{w \in W} BwJ$ such that $\phi_{w,\chi}$ is supported on $Bw^{-1}J$ and $\phi_{w,\chi}|_{w^{-1}J} = 1$. Let $M_w : V(\chi) \rightarrow V({}^w \chi)$ be the intertwining operator defined by

$$(M_w f)(g) = \int_{N \cap wN_w^{-1}} f(w^{-1}ng) \, dn.$$

Then $\{f_w\}$ is the dual basis of the linear functionals $V(\chi) \rightarrow \mathbf{C}$, $f \mapsto (M_w f)(1)$, $w \in W$. Casselman [5] asks for the transition matrix between these two bases, which is in general a very difficult problem. It is better to use the basis

$$\psi_{x,\chi} = \sum_{w \geq x} \phi_{w,\chi}$$

instead of $\phi_{w,\chi}$, and by Möbus inversion one has

$$\phi_{x,\chi} = \sum_{w \geq x} (-1)^{\ell(w)-\ell(x)} \psi_{w,\chi}.$$

If we write $\psi_{x,\chi} = \sum_{w \in W} m(x, w) f_w$, then obviously $m(x, w) = (M_w \psi_{x,\chi})(1)$ and in [4] it is shown that $(m(x, w))$ is upper triangular. In [4] it is conjectured that

$$m(x, w) = \prod_{\alpha \in S(x, w)} \frac{1 - q^{-1} \mathbf{z}^\alpha}{1 - \mathbf{z}^\alpha}$$

when the root system Φ is simply-laced and $|S(x, w)| = \ell(w) - \ell(x)$, and it is proved under the additional assumption that w admits a good word for x .

Let H be the Iwahori–Hecke algebra, which consists of bi- J -invariant functions supported on K . Then H has a basis $\{t_w | w \in W\}$, where t_w is the characteristic function of JwJ , and H is generated by $t_i := t_{\sigma_i}$, $i \in I$. Let $\alpha_\chi : V(\chi)^J \rightarrow H$ be the isomorphism of left H -modules defined by $(\alpha_\chi f)(g) = f(g^{-1})|_K$. Let $\mathcal{M}_w = \mathcal{M}_{w,\mathbf{z}} : H \rightarrow H$ be the map making the following diagram commute:

$$\begin{CD} V(\chi) @>M_w>> V({}^w\chi) \\ @V\alpha_\chi VV @VV\alpha_w\chi V \\ H @>\mathcal{M}_w>> H \end{CD}$$

Define $\mu_{\mathbf{z}}(w) = \mathcal{M}_w(1_H) \in H$. Then

$$\mu_{\mathbf{z}}(\sigma_i) = q^{-1} t_i + (1 - q^{-1}) \frac{\mathbf{z}^{a_i}}{1 - \mathbf{z}^{a_i}}, \tag{6.1}$$

and for $\ell(w_1 w_2) = \ell(w_1) + \ell(w_2)$ one has

$$\mu_{\mathbf{z}}(w_1 w_2) = \mu_{\mathbf{z}}(w_2) \mu_{w_2 \mathbf{z}}(w_1). \tag{6.2}$$

Define $\psi(x) = \alpha_\chi(\psi_x) \in H$. Then $\psi(x) = \sum_{w \geq x} t_w$ is independent of χ . For $f \in H$ let $\Lambda(f)$ be the coefficient of 1 in the expression of f in terms of the basis t_w . Then

$$m(x, w) = \Lambda(\psi(x) \mu_{\mathbf{z}}(w)).$$

For $f, g \in H$ and $x \in W$, write $f - g \geq x$ if $f - g$ is a linear combination of t_w 's with $w \geq x$.

Proposition 6.1 [4]. *Let $s = s_\alpha \in S$, $x \in W$ such that $xs > x$. Then*

$$\psi(x) \mu_{\mathbf{z}}(s) = \frac{1 - q^{-1} \mathbf{z}^\alpha}{1 - \mathbf{z}^\alpha} \psi(x), \quad \psi(xs) \mu_{\mathbf{z}}(s) - \psi(x) \geq xs.$$

Now we can give our formula for $m(x, w)$ in full root system generality, assuming that Condition (B) holds.

Theorem 6.2. *Assume Condition (B). Let $\gamma_k = s_n \cdots s_{i_k+1} \alpha_{i_k}$, $k = 1, \dots, d$. Then*

$$m(x, w) = \prod_{k=1}^d \frac{1 - q^{-1} \mathbf{z}^{\gamma_k}}{1 - \mathbf{z}^{\gamma_k}}.$$

Proof. The proof follows the argument in [4], but we shall give some details for the sake of completeness. Write $\mu(s_n) = \mu_{\mathbf{z}}(s_n)$, $\mu(s_{n-1}) = \mu_{s_n(\mathbf{z})}(s_{n-1})$, ..., suppressing the dependence of spectral parameters. Write $\psi(x)\mu_{\mathbf{z}}(w)$ as a sum

$$\begin{aligned} & [\psi(s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_d} \cdots s_n) \mu(s_n) - \psi(s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_d} \cdots s_{n-1})] \mu(s_{n-1}) \cdots \mu(s_1) \\ & + [\psi(s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_d} \cdots s_{n-1}) \mu(s_{n-1}) - \psi(s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_d} \cdots s_{n-2})] \mu(s_{n-2}) \cdots \mu(s_1) \\ & + \cdots \\ & [\psi(s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_d}) \mu(s_{i_d}) - C(d) \psi(s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_d})] \mu(s_{i_d-1}) \cdots \mu(s_1) \\ & + C(d) [\psi(s_1 \cdots \hat{s}_{i_1} \cdots s_{i_d-1}) \mu(s_{i_d-1}) - \psi(s_1 \cdots \hat{s}_{i_1} \cdots s_{i_d-2})] \mu(s_{i_d-2}) \cdots \mu(s_1) \\ & + \cdots \\ & C(d) \cdots C(1) [\psi(s_1) \mu(s_1) - \psi(1)] \\ & + C(d) \cdots C(1) \psi(1), \end{aligned}$$

where

$$C(k) = \frac{1 - q^{-1} \mathbf{z}^{\gamma_k}}{1 - \mathbf{z}^{\gamma_k}}, \quad k = 1, \dots, d.$$

We will show that the linear functional Λ annihilates every summand except the last, so that $m(x, w) = C(d) \cdots C(1)$.

Since we have the reduced words $w_k^+ := s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_k} s_{i_k+1} \cdots s_n$, $k = 1, \dots, d$, which form the maximal chain $\mathcal{C}_{x, w}^+$, we see that $s_1 \cdots \hat{s}_{i_1} \cdots s_{i_k} > s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_k}$. Therefore, by Proposition 6.1, the summands of the form

$$\prod_{j>k} C(j) [\psi(s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_k}) \mu(s_{i_k}) - C(k) \psi(s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_k})] \mu(s_{i_k-1}) \cdots \mu(s_1)$$

are all equal to zero. Note that the spectral parameter of $\mu(s_{i_k})$ is $s_{i_k+1} \cdots s_n \mathbf{z}$ and one has $(s_{i_k+1} \cdots s_n \mathbf{z})^{\alpha_{i_k}} = \mathbf{z}^{s_n \cdots s_{i_k+1} \alpha_{i_k}} = \mathbf{z}^{\gamma_k}$.

Every other summand is a constant multiple of the form

$$[\psi(s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_2} \cdots s_j) \mu(s_j) - \psi(s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_2} \cdots s_{j-1})] \mu(s_{j-1}) \cdots \mu(s_1). \tag{6.3}$$

Since $s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_2} \cdots s_j$ is reduced, by Proposition 6.1 we have

$$\psi(s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_2} \cdots s_j) \mu(s_j) - \psi(s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_2} \cdots s_{j-1}) \geq s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_2} \cdots s_j.$$

Applying (6.1), (6.2) and arguing as in [4] one can deduce that (6.3) is annihilated by Λ unless

$$s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_2} \cdots s_j \leq s_1 \cdots s_{j-1} = s_1 \cdots s_{j-1} \hat{s}_j. \tag{6.4}$$

Assume that (6.4) is true; let $d' = \max\{1 \leq k \leq d : i_k < j\}$, $x' = s_1 \cdots \hat{s}_{i_1} \cdots \hat{s}_{i_{d'}} \cdots s_j$ and $\mathbf{w}' = s_1 \cdots s_j$. Recall that we have the reduced words $w_k^- := s_1 \cdots \hat{s}_{i_{d-k+1}} \cdots \hat{s}_{i_d} \cdots s_n$, $k = 1, \dots, d$, which make the maximal chain $\mathcal{C}_{x, \mathbf{w}}^-$. Consider the following subchain of $\mathcal{C}_{x, \mathbf{w}}^-$

$$w_{d-d'}^- \rightarrow w_{d-d'+1}^- \rightarrow \cdots \rightarrow w_d^- = x.$$

By taking reduced subwords, it gives rise to a maximal chain of $[x', \mathbf{w}']$

$$\mathcal{C} : \mathbf{w}' = w'_0 \rightarrow w'_1 \rightarrow \cdots \rightarrow w'_{d'} = x'$$

where $w'_i = s_1 \cdots \hat{s}_{d'-i+1} \cdots \hat{s}_{d'} \cdots s_j$, $i = 1, \dots, d'$. Then $\lambda(\mathcal{C}) = (i_{d'}, \dots, i_1)$ is decreasing, which implies that $\mathcal{C} = \mathcal{C}_{x', \mathbf{w}'}^-$. But similarly to the proof of Lemma 5.3(i), this contradicts (6.4) because $\mathcal{C}_{x', \mathbf{w}'}^-$ is lexicographically maximal. This finishes the proof of the theorem. □

Remark 6.3. Given the equivalence of Conditions (A) and (B), proved in the Appendix, Theorem 6.2 gives a correction to the Bump and Nakasuji result [4] in full root system generality. In fact [4, Theorem 1.9] does not hold for the two examples in Remark 5.1. In these cases w admits a good word \mathbf{w} for x but $m(x, w)$ does not factor like (1.3) in [4]. Actually in these cases the Schubert variety $X(w_0x)$ is singular at e_{w_0w} (in the notations of [1]), so $m(x, w)$ never factors. (It is also easy to check by computer that $m(x, w)$ does not factor in these cases.) Moreover, the type D_4 example in Remark 5.1 also gives a counterexample to Conjecture 1.2 (and (1.3)) in paper [4]. In this case the root system is simply-laced and $|S(x, w)| = \ell(w) - \ell(x)$, but the Kazhdan–Lusztig polynomial is $P_{w_0w, w_0x} = (1 + q)^2$, which implies singularity and $m(x, w)$ never factors as stated in [4, Conjecture 1.2].

Appendix A. Proof of Conjecture 5.9

BY DINAKAR MUTHIAH¹ AND ANNA PUSKÁS²

In this appendix, we prove Conjecture 5.9. The conjecture is that the following three conditions are equivalent.

- (i) $\lambda_{x, \mathbf{w}} = \lambda(\mathcal{C}_{x, \mathbf{w}}^-)^*$;
- (ii) $\lambda(\mathcal{C}_{x, \mathbf{w}}^+) = \lambda(\mathcal{C}_{x, \mathbf{w}}^-)^*$;
- (iii) $\lambda_{x, \mathbf{w}} = \lambda(\mathcal{C}_{x, \mathbf{w}}^+)$.

As mentioned right below Conjecture 5.9, it suffices to prove that (i) and (ii) are equivalent.

We will keep the notations in the previous sections. In particular, we have two elements $x \leq w$ of W , $\mathbf{w} = s_1 \cdots s_n$ a reduced word for w ; $\lambda_{x, \mathbf{w}} = (\lambda_1, \dots, \lambda_k)$, $\lambda(\mathcal{C}_{x, \mathbf{w}}^+) = (i_1, \dots, i_d)$ and $\lambda(\mathcal{C}_{x, \mathbf{w}}^-)^* = (j_1, \dots, j_d)$.

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A.1. Proof of (i) \implies (ii)

Lemma A.1. *Let x, w, \mathfrak{w} be as before, $\lambda_{x,\mathfrak{w}} = (\lambda_1, \dots, \lambda_k)$, and $\lambda(\mathcal{C}_{x,\mathfrak{w}}^+) = (i_1, \dots, i_d)$. Then $i_1 = 1$ if and only if $\lambda_1 = 1$.*

Proof. Assume first that $i_1 = 1$. Then the chain $\mathcal{C}_{x,\mathfrak{w}}^+$ starts with $\mathfrak{w}\bar{\gamma}$, and hence $x \leq \mathfrak{w}\bar{\gamma}$ and thus $\lambda_1 = 1$. For the other direction, assume $\lambda_1 = 1$. Omitting the first simple reflection from \mathfrak{w} only decreases its length by 1; hence $\ell(\mathfrak{w}\bar{\gamma}) = \ell(w) - 1$. Composing $w \rightarrow \mathfrak{w}\bar{\gamma}$ with a maximal chain from $\mathfrak{w}\bar{\gamma}$ to x gives a maximal chain \mathcal{C} from \mathfrak{w} to x whose label starts with 1. Then $\mathcal{C}_{x,\mathfrak{w}}^+ \leq_L \mathcal{C}$ implies $i_1 = 1$. \square

Remark A.2. If (i) holds for x and \mathfrak{w} , i.e., $\lambda_{x,\mathfrak{w}} = \lambda(\mathcal{C}_{x,\mathfrak{w}}^-)^*$, then \mathfrak{w} is a good word of w for x . (Omitting all the reflections from \mathfrak{w} that appear in $\lambda_{x,\mathfrak{w}}$ is the same as taking the last element of the maximal chain $\mathcal{C}_{x,\mathfrak{w}}^-$; that last element is x .)

Proposition A.3. (i) \implies (ii).

Proof. We proceed by induction on $\ell(w) + (\ell(w) - \ell(x))$; the base case is trivial.

Assume that (i) holds for a pair x, \mathfrak{w} , i.e., $\lambda_{x,\mathfrak{w}} = \lambda(\mathcal{C}_{x,\mathfrak{w}}^-)^*$. We would like to show that (ii) holds for x, \mathfrak{w} as well, i.e., $\lambda(\mathcal{C}_{x,\mathfrak{w}}^+) = \lambda(\mathcal{C}_{x,\mathfrak{w}}^-)^*$.

Consider $\lambda_1 = j_1$, the first index in the labels $\lambda_{x,\mathfrak{w}} = \lambda(\mathcal{C}_{x,\mathfrak{w}}^-)^*$. We distinguish between two cases according to whether $\lambda_1 = j_1 = 1$ or $\lambda_1 = j_1 > 1$.

Case 1: $\lambda_1 = j_1 = 1$. Then by Lemma 5.4(iv), we have that (i) holds for the pair $x, \mathfrak{w}' = s_1\mathfrak{w}$. Then by induction, (ii) holds for x and \mathfrak{w}' , i.e., $\lambda(\mathcal{C}_{x,\mathfrak{w}'}^+) = \lambda(\mathcal{C}_{x,\mathfrak{w}'}^-)^*$.

By Lemma 5.3(iii) and (iv), we have:

$$(i_2 - 1, \dots, i_d - 1) = (j_2 - 1, \dots, j_d - 1). \tag{A 1}$$

Together with $i_1 = 1$ (Lemma A.1) we conclude that $i_r = j_r$ for every $1 \leq r \leq d$; hence (ii) holds for the pair x, \mathfrak{w} .

Case 2: $\lambda_1 = j_1 > 1$. By Lemmas 5.2(iii) and 5.3(ii), (i) holds for the pair $x' = s_1x$ and $\mathfrak{w}' = s_1\mathfrak{w}$. By induction, (ii) also holds for x', \mathfrak{w}' . By Lemma A.1, $i_1 > 1$. Thus by Lemma 5.3(i) and (ii), we have $(i_1 - 1, \dots, i_d - 1) = (j_1 - 1, \dots, j_d - 1)$, which implies $(i_1, \dots, i_d) = (j_1, \dots, j_d)$. \square

A.2. Proof of (ii) \implies (i)

Lemma A.4. *Let x, w, \mathfrak{w} be as before. Write $\lambda(\mathcal{C}_{x,\mathfrak{w}}^+) = (i_1, \dots, i_d)$, $\lambda(\mathcal{C}_{x,\mathfrak{w}}^-)^* = (j_1, \dots, j_d)$.*

Suppose $j_1 = 1$; then:

- $x \leq s_1w$;
- $\lambda_{x,s_1\mathfrak{w}} = (\lambda_{x,\mathfrak{w}} \setminus \{1\}) - 1$.

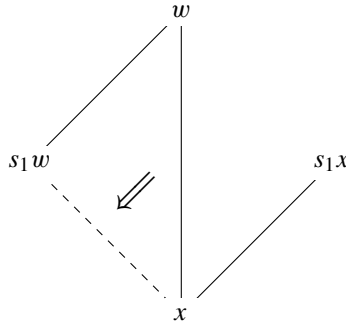
Suppose $i_1 > 1$; then:

- $s_1x \leq s_1w$;
- $\lambda_{s_1x,s_1\mathfrak{w}} \supseteq \lambda_{x,\mathfrak{w}} - 1$.

Here we write $(\lambda_{x,w} \setminus \{1\}) - 1$ and $\lambda_{x,w} - 1$ to refer to the set obtained by subtracting 1 from all elements.

Proof. Note that $s_1 w$ is a reduced word for $s_1 w$, and $s_1 w < w$.

First suppose $j_1 = 1$. By Lemma 5.3(iv) we have $x < s_1 x$. By Lemma 3.1 we may draw the diagram

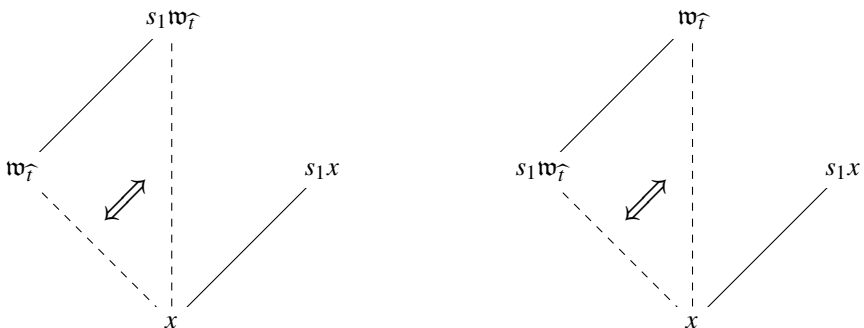


and conclude that $x \leq s_1 w$. Let $1 < t \leq n$ and $w_{\widehat{t}} := s_1 \cdots \widehat{s_t} \cdots s_n$, and $s_1 w_{\widehat{t}} := s_2 \cdots \widehat{s_t} \cdots s_n (= (s_1 w)_{\widehat{t-1}})$. To show $\lambda_{x,s_1 w} = (\lambda_{x,w} \setminus \{1\}) - 1$, it suffices to prove

$$x \leq w_{\widehat{t}} \iff x \leq s_1 w_{\widehat{t}}. \tag{A 2}$$

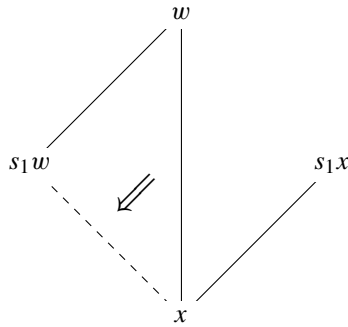
(Note that we are slightly abusing notation. For example, when we write $x \leq w_{\widehat{t}}$, we mean $x \leq w_{\widehat{t}}$, where $w_{\widehat{t}}$ is the Weyl group element obtained by multiplying out the word $w_{\widehat{t}}$.)

To prove (A 2), we use Lemma 3.1 again. We have either $w_{\widehat{t}} < s_1 w_{\widehat{t}}$ or $w_{\widehat{t}} > s_1 w_{\widehat{t}}$; we may accordingly draw one of the following two diagrams.

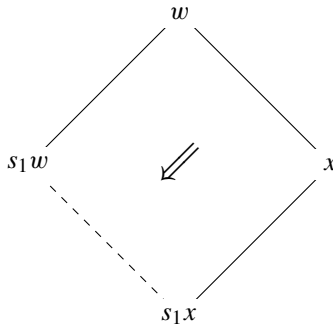


These diagrams together imply that (A 2) holds in both cases.

Next suppose $i_1 > 1$. The argument in this case is very similar to the one above. We claim $x > s_1 x$. Assume to the contrary that $x < s_1 x$. Then again by Lemma 3.1 we may draw the following diagram.



This contradicts the statement of Lemma 5.3(i) that $x \not\leq s_1w$. Hence we have $x > s_1x$, and consequently the diagram



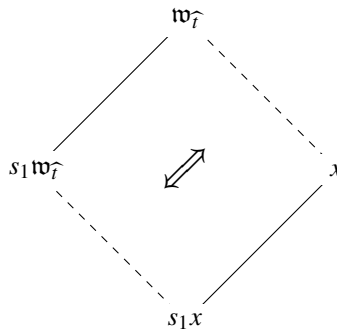
shows that $s_1x \leq s_1w$. Take $1 < t \leq n$ and w_t and s_1w_t as in the case $i_1 = 1$ above. To prove $\lambda_{s_1x, s_1w} \geq \lambda_{x, w} - 1$ we need to show

$$x \leq w_t \implies s_1x \leq s_1w_t. \tag{A 3}$$

First consider the case when $w_t < s_1w_t$. Then if $x \leq w_t$ we have

$$s_1x < x \leq w_t < s_1w_t, \tag{A 4}$$

whence $s_1x < s_1w_t$. If on the other hand $w_t > s_1w_t$, then again by Lemma 3.1 we have the diagram



which proves (A 3). □

Proposition A.5. (ii) \implies (i).

Proof. Let x, w, \mathfrak{w} be as before. We proceed by induction on $\ell(w) + (\ell(w) - \ell(x))$; the base case is trivial.

Let us assume $\lambda(\mathcal{C}_{x,\mathfrak{w}}^+) = \lambda(\mathcal{C}_{x,\mathfrak{w}}^-)^*$. Write $\lambda_{x,\mathfrak{w}} = (\lambda_1, \dots, \lambda_k)$, $\lambda(\mathcal{C}_{x,\mathfrak{w}}^+) = (i_1, \dots, i_d)$, and $\lambda(\mathcal{C}_{x,\mathfrak{w}}^-) = (j_d, \dots, j_1)$. Our assumption means that $i_r = j_r$ for all r .

Case 1: $i_1 = j_1 = 1$. In this case $\lambda_1 = 1$ by Lemma A.1, and Lemma A.4 tells us that $x \leq s_1 w$ and

$$\lambda_{x,s_1\mathfrak{w}} = (\lambda_{x,\mathfrak{w}} \setminus \{1\}) - 1. \tag{A 5}$$

Then by Lemma 5.3(iii) and (iv), we have that:

$$\lambda(\mathcal{C}_{x,s_1\mathfrak{w}}^+) = \lambda(\mathcal{C}_{x,s_1\mathfrak{w}}^-)^* = (i_2 - 1, \dots, i_d - 1). \tag{A 6}$$

By induction, we know:

$$\lambda_{x,s_1\mathfrak{w}} = \lambda(\mathcal{C}_{x,s_1\mathfrak{w}}^-)^*. \tag{A 7}$$

By (A 5)–(A 7), $\lambda_{x,\mathfrak{w}} = (i_1, \dots, i_d)$. Therefore $\lambda_{x,\mathfrak{w}} = \lambda(\mathcal{C}_{x,\mathfrak{w}}^-)^*$.

Case 2: $i_1 = j_1 > 1$. The argument is very similar. In this case, $\lambda_1 > 1$ by Lemma A.1, and Lemma A.4 tells us that $s_1 x \leq s_1 w$ and

$$\lambda_{s_1 x, s_1 \mathfrak{w}} \supset \lambda_{x,\mathfrak{w}} - 1. \tag{A 8}$$

By Lemma 5.3(i) and (ii), we have that:

$$\lambda(\mathcal{C}_{s_1 x, s_1 \mathfrak{w}}^+) = \lambda(\mathcal{C}_{s_1 x, s_1 \mathfrak{w}}^-)^* = (i_1 - 1, \dots, i_d - 1). \tag{A 9}$$

By induction, we know:

$$\lambda_{s_1 x, s_1 \mathfrak{w}} = \lambda(\mathcal{C}_{s_1 x, s_1 \mathfrak{w}}^-)^*. \tag{A 10}$$

In particular:

$$\#\lambda_{s_1 x, s_1 \mathfrak{w}} = \ell(w) - \ell(x). \tag{A 11}$$

By Deodhar’s inequality:

$$\#\lambda_{x,\mathfrak{w}} \geq \ell(w) - \ell(x). \tag{A 12}$$

So (A 8), (A 11), and (A 12) together imply:

$$\lambda_{s_1 x, s_1 \mathfrak{w}} = \lambda_{x,\mathfrak{w}} - 1. \tag{A 13}$$

By (A 9), (A 10) and (A 13), $\lambda_{x,\mathfrak{w}} = (i_1, \dots, i_d)$. Therefore $\lambda_{x,\mathfrak{w}} = \lambda(\mathcal{C}_{x,\mathfrak{w}}^-)^*$. □

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